

## Lecture 19

### AoI and Sampling

Reading: Wait or Update TIT 2017.

JSAC AoI survey

Sun, Cyr 2019.

Ornee, Sun 2020



$p(t)$  is non-decreasing.

$$\bar{P}_{opt} = \inf_{\pi \in \Pi} \limsup_{T \rightarrow \infty} \frac{1}{T} E \left[ \int_0^T P(\Delta(t)) dt \right].$$

$\pi: (s_1, s_2, \dots)$  is a sampling policies  
↓  
Sampling time.

$\Pi$ : the set of causal policies.

---

Theorem:

The policy  $(s_1(\beta), s_2(\beta), \dots)$  defined by

$$s_{i+1}(\beta) = \inf \left\{ t \geq D_i(\beta) : E [ P(\Delta(t + Y_{i+1})) ] \geq \beta \right\},$$

where  $D_i(\beta) = s_i(\beta) + Y_i$ ,  $\Delta(t) = t - s_i(\beta)$ ,

and  $\beta$  is the root of

$$E \left[ \int_{D_i(\beta)}^{D_{i+1}(\beta)} P(\Delta(t)) dt \right] = \beta E [ D_{i+1}(\beta) - D_i(\beta) ]. \quad (*)$$

is an optimal solution to (41).

Further  $\beta = \bar{P}_{opt}$ .

Sample  $i+1$  is generated at the earliest time  $t$  satisfying:

(i)  $t \geq D_i(\beta)$ , sample  $i$  has been delivered.

(ii)  $E[p(\Delta(t+Y_{i+1}))]$  has grown to  $\geq$  threshold  $\beta = \bar{P}_{opt}$ .

Note: (\*) has a unique solution.

Reading: Theorem 1  
Sun, Cyr 2019.

Theorem 4.

Ornee, Sun 2020

Proof.

Step 1.

Lemma 1. It is sub-optimal to take a new sample before the previous sample is delivered.

Proof. If a sample is taken when the server is busy, one can design a better sampler, by postponing the sampling time to a later time when the server becomes idle.  $\square$

By this lemma, we only need to consider the policies

$$\Pi_1 = \{ \pi \in \Pi : S_{i+1} \geq D_i = S_i + Y_i \text{ for all } i \}.$$

Let  $Z_i = S_{i+1} - D_i \geq 0$ , be the waiting time between  $D_i$  and  $S_{i+1}$ .

Designing the sampling times  $(S_1, S_2, \dots)$  is equivalent to designing the waiting times  $(Z_1, Z_2, \dots)$ .

Step 2.

Suppose that  $\forall \pi \in \Pi_1$ .

$$(a) \limsup_{T \rightarrow \infty} \frac{1}{T} E \left[ \int_0^T P(\Delta(t)) dt \right] = \lim_{i \rightarrow \infty} \frac{E \left[ \int_0^{D_i} P(\Delta(t)) dt \right]}{E[D_i]}.$$

$$(b) \lim_{i \rightarrow \infty} \frac{1}{i} E[D_i] \quad \lim_{i \rightarrow \infty} \frac{1}{i} E \left[ \int_0^{D_i} P(\Delta(t)) dt \right] \text{ exist.}$$

The assumptions will be discussed later.

Then, we want to minimize.

$$\limsup_{T \rightarrow \infty} \frac{1}{T} E \left[ \int_0^T P(\Delta(t)) dt \right]$$

$$= \lim_{i \rightarrow \infty} \frac{E \left[ \int_0^{D_i} P(\Delta(t)) dt \right]}{E[D_i]}.$$

$$= \lim_{i \rightarrow \infty} \frac{\sum_{j=1}^i E \left[ \int_{D_i}^{D_{i+1}} P(\Delta(t)) dt \right]}{\sum_{j=1}^i E[Y_j + Z_j]}$$

$$\begin{cases} D_i = S_i + Y_i \\ S_{i+1} = D_i + Z_i \\ D_{i+1} = S_{i+1} + Y_{i+1}. \end{cases}$$

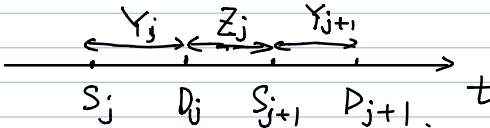
$$\Delta(t) = t - \max \{ S_i : D_i \leq t \}.$$

because  $D_i \leq D_{i+1} \quad \forall i$ , then

$$\Delta(t) = t - S_i \quad \text{if } D_i \leq t < D_{i+1}.$$

$$\int_{D_j}^{D_{j+1}} P(\Delta(t)) dt$$

$$= \int_{D_j}^{D_{j+1}} P(t - S_j) dt.$$



$$\tau = t - S_j.$$

$$= \int_{Y_j}^{Y_j + Z_j + Y_{j+1}} P(\tau) d\tau$$

if  $t = D_j$ , then

$$\tau = t - S_j = Y_j.$$

if  $t = D_{j+1}$ , then

$$\tau = t - S_j$$

$$= D_{j+1} - S_j$$

$$= Y_j + Z_j + Y_{j+1}.$$

Def:  $q(Y, Z, Y')$

$$= \int_Y^{Y+Z+Y'} P(\tau) d\tau.$$

The problem is reformulated as

$$\bar{P}_{opt} = \inf_{\pi \in \Pi_1} \lim_{i \rightarrow \infty} \frac{\sum_{j=1}^i E [q(Y_j, Z_j, Y_{j+1})]}{\sum_{j=1}^i E (Y_j + Z_j)}. \quad (1)$$

Step 3.

Consider the following problem:

$$h(c) = \inf_{\pi \in \Pi_1} \lim_{i \rightarrow \infty} \frac{1}{i} \sum_{j=1}^i E [q(Y_j, Z_j, Y_{j+1}) - c(Y_j + Z_j)]. \quad (2)$$

Lemma. (a)  $\bar{P}_{opt} \begin{matrix} \geq \\ \leq \end{matrix} c$  if & only if  $h(c) \begin{matrix} \geq \\ < \end{matrix} 0$

(b) If  $h(c) = 0$ , then the solutions to (1) and (2) are identical.

Proof: "If  $\bar{P}_{opt} \leq c$ , then  $h(c) \leq 0$ ".

If  $\bar{P}_{opt} \leq c$ , then for any  $\varepsilon > 0$ .

there exists a policy  $\pi = (z_1, z_2, \dots)$  satisfying

$$\lim_{i \rightarrow \infty} \frac{\sum_{j=1}^i E [q(Y_j, z_j, Y_{j+1})]}{\sum_{j=1}^i E (Y_j + z_j)} \leq c + \varepsilon.$$

$$\lim_{i \rightarrow \infty} \frac{\sum_{j=1}^i E [q(Y_j, z_j, Y_{j+1})] - c \sum_{j=1}^i E [Y_j + z_j]}{\sum_{j=1}^i E (Y_j + z_j)} \leq \varepsilon.$$

$$E [Y_j] > 0.$$

$$\frac{1}{i} \sum_{j=1}^i E (Y_j + z_j) \geq m > 0.$$

$$\lim_{i \rightarrow \infty} \frac{1}{i} \sum_{j=1}^i E [Y_j + z_j] \text{ exist.}$$

$$\begin{aligned} & \lim_{i \rightarrow \infty} \frac{1}{i} \sum_{j=1}^i E [q(Y_j, z_j, Y_{j+1})] - c \frac{1}{i} \sum_{j=1}^i E [Y_j + z_j] \\ & \leq \varepsilon \frac{1}{i} \sum_{j=1}^i E [Y_j + z_j] \end{aligned}$$



if  $\lim_{i \rightarrow \infty} a_i = a$ ,  $\lim_{i \rightarrow \infty} b_i = b$ , then

$$\lim_{i \rightarrow \infty} \frac{a_i}{b_i} = \frac{a}{b}, \quad \lim_{i \rightarrow \infty} a_i b_i = a b.$$

The choice of  $\varepsilon > 0$ , is arbitrary.

$$\inf \left\{ \lim_{i \rightarrow \infty} \frac{1}{i} \sum_{j=1}^i E [q_r(Y_j, Z_j, Y_{j+1})] - c \frac{1}{i} \sum_{j=1}^i E [Y_j + Z_j] \right\} \leq 0.$$

$$h(c) \leq 0.$$

---

"If  $h(c) \leq 0$ , then  $\bar{P}_{opt} \leq c$ ".

Next we need

$$\begin{aligned} \bar{P}_{opt} < c & \text{ if \& only if } h(c) < 0. \\ \bar{P}_{opt} \geq c & \text{ if \& only if } h(c) \geq 0. \end{aligned}$$

By this,

" $\bar{P}_{opt} \leq c$  if \& only if  $h(c) \leq 0$ ".